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- Statement Regarding Fed sponsored R & D							
- Reference to Microfiche Appendix		c.				g identity of above copie	98
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- Brief Description of the Drawings (if filed)			•	•	•	r sheet & document(s))	
- Detailed Description				.R.§3.73(b) <i>there is an a</i>			
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see 37 C.F.R. §§ 1.63(d)(2) and 1.			Other:			9	
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- Specification (26) pages;
- Claims (18) pages;
- Abstract (1) pages;
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1	This applica	ation is submitted in the na	ame of the following inventor(s):
2			
3	Inventor	Citizenship	Residence City and State
4	Scott SCHOENTHAL	United States	San Ramon, California
5			
6	The assigne	ee is Network Appliance, I	Inc., a California corporation having an
7	office at 495 East Java D	rive, Sunnyvale CA 94089	9.
8			
9		TITLE OF THE INV	ENTION
8 9 10 11			
11	Persis	tent and Reliable Delivery	of Event Messages
12			
12 13 14	F	BACKGROUND OF THE	INVENTION
14			
15	1. Field of the Invent	tion	
16			
17	This invent	tion relates to persistent an	nd reliable delivery of event messages
18	including event messages	s in file server systems in	which it is desired to maintain reliable
19	file system consistency.		
20			

#### Related Art 2.

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3 ers and similar devices, it often occurs that the system, or some subsystem within that 4 system, generates a message indicating the occurrence of a special event. Typically, the 5 special event is an error of some kind, and the message conveys information regarding 6 the nature of the special event, such as the type of error and the subsystem within which 7 the error occurred. Many systems that provide services, including file servers, make ef-9 10 11 forts to assure that the services are reliably provided, and that the system providing the services is in a consistent state at all times. Thus, such systems find it advantageous to assure that all state information regarding the system, including state information relating <u>1</u>2 to error messages, is persistently and reliably maintained. Such systems also find it advantageous to assure that all event messages are reliably delivered and persistently main-

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Accordingly, it would be advantageous to provide a technique for persistent and reliable delivery of event messages, that is not subject to the drawbacks of the known art. Preferably, those parts of the system responsible for delivering event messages are able to persistently maintain those event messages until delivery of those event messages has been confirmed by the intended recipient of the event message. Moreover, those parts of the system responsible for recovering from system crashes and other system er-

tained until delivery is confirmed by the intended recipient of the event message.

In systems that provide services to clients, such as those including file serv-

rors are able to persistently maintain those event messages until delivery, even after recovery from system crashes or other system errors.

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## SUMMARY OF THE INVENTION

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The invention provides a method and system for persistent and reliable delivery of event messages, that is not subject to the drawbacks of the known art. Those parts of the system responsible for delivering event messages are able to persistently maintain those event messages until delivery of those event messages has been confirmed by the intended recipient of the event message. Those parts of the system responsible for recovering from system crashes and other system errors are able to persistently maintain those event messages until delivery, even after recovery from system crashes or other system errors.

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In a first aspect of the invention, the system includes a set of event message producers, and maintains an event-indication queue of those event messages provided by the event producers using a set of pre-allocated resources. This first aspect allows the system to maintain the event-indication queue even when the event message indicates that allocation of new resources (such as for recording event messages) is unstable. An event-distribution engine distributes event messages to intended recipients and, after having received confirmation that the event messages were received, removes them from the event-indication queue. Recipients (also called "consumers") of event messages re-

ceive the event messages, acknowledge their receipt thereof, and might take action in response to the event message.

In a second aspect of the invention, the system includes a persistent memory (such as NVRAM or other non-volatile memory), in which event messages can be recorded until they are completely handled by the event-distribution engine. Upon recovery from a system crash or other system error, a replay-event producer retrieves those event messages recorded in the persistent memory and not yet completely handled, and re-presents them as event messages for the event-indication queue and the event-distribution engine.

In a third aspect of the invention, the system includes an initialization memory (also called a "system boot memory"), in which event messages can be recorded until the system has completed its operations of initializing and becoming completely prepared to handle event messages. Upon recovery from a system crash or other system error, the replay-event producer retrieves those event messages recorded in the initialization memory and re-presents them as event messages for the event-indication queue and the event-distribution engine.

In a fourth aspect of the invention, a cluster of file servers collectively forming a highly-available system shares persistent memories. Each individual file server writes event messages to both its own and at least one other persistent memory, so

that upon a system crash or other system error, at least one other file server has a record of those event messages that were presented by event producers but not yet completely handled by the event-distribution engine. Upon indication of a system crash or other system error by a first file server in the cluster, a second file server in the cluster uses its replay-event producer. Thus, the second file server retrieves those event messages recorded in the persistent memory and not yet completely handled by the first file server, and re-presents them as event messages for its own event-indication queue and eventdistribution engine.

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In a fifth aspect of the invention, a cluster of event recipients can be coupled to a single multiplexing recipient that includes a second persistent memory in which event messages are recorded after receipt and before being redistributed or otherwise completely handled. Upon recovery from a system crash or other system error by the multiplexing recipient, the multiplexing recipient includes a replay-event producer that retrieves those event messages recorded in the second persistent memory, and re-presents

them as event messages as if newly received from an event producer.

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The invention provides an enabling technology for a wide variety of applications for persistent and reliable delivery of event messages, so as to obtain substantial advantages and capabilities that are novel and non-obvious in view of the known art. Examples described below primarily relate to reliable file systems, but the invention is

1	broadly applicable to many different types of systems in which persistent and reliable de-
2	livery of event messages is desired.

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## BRIEF DESCRIPTION OF THE DRAWINGS

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Figure 1 shows a block diagram of a portion of a system capable of persis-6 tent and reliable delivery of event messages. 7

Figure 2 shows a process flow diagram of a method for operating a system as in figure 1.

### DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

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In the following description, a preferred embodiment of the invention is described with regard to preferred process steps and data structures. Embodiments of the invention can be implemented using general-purpose processors or special purpose processors operating under program control, or other circuits, adapted to particular process steps and data structures described herein. Implementation of the process steps and data structures described herein would not require undue experimentation or further invention.

1	Relate	ed Applications
2		
3		Inventions described herein can be used in conjunction with technology de-
4	scribe	ed in the following documents.
5		
6	•	U.S. Patent Application Serial No, Express Mail Mailing No. EL
7		524781089US, filed August 18, 2000, in the name of Blake LEWIS, attorney
_ 8		docket number 103.1033.01, titled "Reserving File System Blocks"
道 河 9		
8 9 0	•	U.S. Patent Application Serial No, Express Mail Mailing No.
11		EL524780242US, filed August 18, 2000, in the name of Rajesh SUNDARAM,
12		attorney docket number 103.1034.01, titled "Dynamic Data Storage"
12 13 14		
<u>-</u> 14	•	U.S. Patent Application Serial No, Express Mail Mailing No.
15		EL524780239US, filed August 18, 2000, in the name of Blake LEWIS, attorney
16		docket number 103.1035.01, titled "Instant Snapshot"
17		
18	•	U.S. Patent Application Serial No, Express Mail Mailing No.
19		EL524781092US, filed August 18, 2000, in the name of Douglas P. DOUCETTE,
20		attorney docket number 103.1045.01, titled "Improved Space Allocation in a
21		Write Anywhere File System"

1	and
2	• U.S. Patent Application Serial No, Express Mail Mailing No.
3	EL524780256US, filed August 18, 2000, in the name of Ray CHEN, attorney
4	docket number 103.1047.01, titled "manipulation of Zombie Files and Evil-Twin
5	Files"
6	
7	Each of these documents is hereby incorporated by reference as if fully set
8	forth herein. This application claims priority of each of these documents. These docu-
8	ments are collectively referred to as the "Incorporated Disclosures."
10 111	
<b>1</b> 1	Lexicography
12	
12 13 14	The following terms refer or relate to aspects of the invention as described
14	below. The descriptions of general meanings of these terms are not intended to be limit-
15	ing, only illustrative.
16	
17	• event messages In general, an event message refers to an alert or notification of
18	a system event, for which an intended recipient of that event message may wish to
19	receive and possibly take action in response to. Examples of event messages in-
20	clude notification of errors, or of events to be logged or otherwise administratively
21	monitored.

• event replay --- In general, re-presenting event messages (from a memory or other record) by a system element other than the one that generated the original event message. Event messages that are replayed are treated substantially identically to originally generated event messages.

• persistent, reliable --- In general, persistent refers to memory or another record that is capable of surviving a disruption such as a system crash or a system error. In general, reliable refers to a process that is capable of being performed atomically or otherwise completely even in the event of a disruption such as a system crash or a system error.

• pre-allocated resources --- In general, resources that have been allocated ahead of time for use by a system element (such as for delivery of event messages), so that system element can operate even when disruptions in system operation make it uncertain that resource allocation will be effective at the time when the system element needs those resources.

• system crashes, system errors --- In general, a system crash or a system error refers to a disruption in system operation sufficiently serious as to place the continuation of system operations (such as delivery of event messages) in doubt. In the description herein, it is assumed that continuation of system operations does not generally survive such disruptions.

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system initialization --- In general, a time during which a system (such as a file server) is booting or initializing, and so is not necessarily able to perform all necessary operations (such as allocation of resources or delivery of event messages).

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As noted above, these descriptions of general meanings of these terms are not intended to be limiting, only illustrative. Other and further applications of the invention, including extensions of these terms and concepts, would be clear to those of ordinary skill in the art after perusing this application. These other and further applications are part of the scope and spirit of the invention, and would be clear to those of ordinary skill in the art, without further invention or undue experimentation.

System Elements

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Figure 1 shows a block diagram of a portion of a system capable of persistent and reliable delivery of event messages.

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A system 100 includes a set of event producers 110, a set of pre-allocated initialization event message resources 120, a set of pre-allocated post-initialization event message resources 130, a persistent memory 140, an event indication queue 150, an event distribution engine 160, a set of event recipients 170 including at least one multiplexing recipient 171 and a set of intended recipients 172, a second persistent memory 180 at the multiplexing recipient 171, and an event replay engine 190.

The event producers 110 include system elements, such as software modules or hardware circuits, each capable of generating at least one event message 111 for delivery to at least one intended recipient 172. In a preferred embodiment, each event message 111 has a standardized format, including information about the time the event was recognized, the system element that recognized the event, the nature of the event, and any detailed information about the event necessary or desirable for the intended recipient 172 to know.

The pre-allocated initialization event message resources 120 include memory, and possibly other resources, for recording and maintaining event messages for further processing. In a preferred embodiment, the pre-allocated initialization event message resources 120 include resources allocated by the system 100 prior to generation of any new event messages 111. In a preferred embodiment, the event messages 111 are recorded in the pre-allocated initialization event message resources 120 before being processed by the event distribution engine 160.

The pre-allocated post-initialization event message resources 130, similar to the pre-allocated initialization event message resources 120, memory, and possibly other resources, for recording and maintaining event messages for further processing. In

a preferred embodiment, the pre-allocated post-initialization event message resources 130 include resources allocated by the system 100 during an initialization period and prior to generation of any new event messages 111. In a preferred embodiment, event messages 111 that are generated during the system initialization period are recorded in the preallocated post-initialization event message resources 130. After the system initialization period, event messages 111 recorded therein are processed by the event distribution engine 160.

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8 9 1 1 1 12 13 14 15 system crash or system error, the persistent memory 140 will still record those event mes-

or other memory whose contents are expected to survive a system crash or system error. In a preferred embodiment, the persistent memory 140 includes NVRAM also used with the WAFL file system. However, in alternative embodiments, the persistent memory 140 may include any other form of persistent memory, whether NVRAM or not, and whether coordinating with aspects of the WAFL file system or not. Thus, upon recovery from a

The persistent memory 140 includes a memory, such as NVRAM, SRAM,

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The event indication queue 150 includes a memory having a queue of information about event messages 111 (such as the event messages 111 themselves).

sages 111 that were not fully processed before the system crash or system error.

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The event distribution engine 160 includes a system element capable of reading information about event messages 111 from the event indication queue 150 and

capable of delivering at those event messages 111 to intended recipients 172 thereof. In a

preferred embodiment, the event distribution engine 160 includes a software module in

the system 100.

The event recipients 170 (including at least one multiplexing recipient 171 and a set of intended recipients 172) include system elements, possibly at remote devices such as clients for the file server system 100, capable of receiving event messages 111 and deciding whether or not to act in response to those event messages 111. In a preferred embodiment, actions take with regard to event messages 111 can include alerts or notification of selected users (such as a system operator), logging the event messages 111, or maintaining statistics with regard thereto.

The second persistent memory 180 at the multiplexing recipient 171 includes, similar to the persistent memory 140, a memory, such as NVRAM, SRAM, or other memory whose contents are expected to survive a system crash or system error. The multiplexing recipient 171 includes a recipient replay element 181, capable of reading information about event messages 111 from the second persistent memory 180 and capable of replaying those event messages 111 as if newly received by the multiplexing recipient 171 (thus delivering those event messages 111 to the intended recipients 172).

The event replay engine 190 includes a system element capable of reading information about event messages 111 from the persistent memory 140 and capable of

103.1048.01 replaying those event messages 111 as if newly generated. The replay element 190 in-1 2 cludes a system initialization replay sub-element 191, an incomplete event distribution 3 replay sub-element 192, and a cooperating systems replay sub-element 193. 4 Method of Operation 5 6 7 Figure 2 shows a process flow diagram of a method for operating a system 8 as in figure 1. 9 10 11 11 11 11 11 11 11 A method 200 includes a set of flow points and a set of steps. The system

100 performs the method 200. Although the method 200 is described serially, the steps of the method 200 can be performed by separate elements in conjunction or in parallel, whether asynchronously, in a pipelined manner, or otherwise. There is no particular requirement that the method 200 be performed in the same order in which this description lists the steps, except where so indicated.

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As described below, the method 200 includes a set of processes, each of which has a set of tasks operating independently and asynchronously with regard to each other.

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#### Processing Event Messages 1.

1	A first process in the method 200 is described with regard to a flow point
2	210, a flow point 220, and steps there-between. This first process includes a set of three
3	tasks, each of which operates independently and asynchronously with regard to each
4	other.
5	
6	At the flow point 210, the system 100 is ready to receive an event message
7	111.
8 9 0 1	Event Generation
<b>1</b> 11	A first task includes a sequence including a step 211, a step 212, a step 213,
	and a step 214. In a preferred embodiment, the first task in its process includes these
12 13 13	steps being performed in sequence and repeatedly.
15	At the step 211, an event producer 110 generates an event message 111.
16	
17	At the step 212, if the system 100 is in normal operation, the event message
18	111 is recorded in the pre-allocated initialization event message resources 120. If the
19	system 100 is still in its initialization time duration, the event message 111 is recorded in
20	the pre-allocated post-initialization event message resources 130.

At the step 213, the system 100 copies information about the event message 111 to a set of locations in the persistent memory 140. In a preferred embodiment, the persistent memory 140 includes a set of memory sections 141, each persistently maintaining system information; at least one of these memory sections 141 maintains information about event messages 111. Other memory sections 141 maintain information about other aspects of the system 100, such as a consistency state of the file system or a set of incomplete file system requests.

In a preferred embodiment, the memory sections 141 associated with event messages 111 maintain information about those event messages 111 in a FIFO having a head pointer and a tail pointer. FIFOs are known in the art of computer data storage. When information about a new event message 111 is recorded in the persistent memory 140, the FIFO is updated to add the information about the new event message 111 to an end of the list. When confirmation is received that the event message 111 was delivered to its intended recipients 172, the FIFO is updated to remove the information about the

event message 111.

At the step 214, similar to the step 213, the system 100 copies information about the event message 111 to the event indication queue 150. In a preferred embodiment, the event indication queue 150 includes a FIFO similar to that maintained in the persistent memory 140.

### Event Distribution

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A second task includes a sequence including a step 215. In a preferred embodiment, the second task in its process includes this step being performed repeatedly.

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At the step 215, the event distribution engine 160 responds to the information about the event message 111 in the event indication queue 150. The event distribution engine 160 delivers the event message 111 to its intended recipients 172. As part of this step, each particular intended recipient 172, when it receives the event message 111, responds to the event distribution engine 160 to confirm its receipt of the event message 111.

## Event Confirmation

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A third task includes a sequence including a step 216 and a step 217. In a preferred embodiment, the third task in its process includes these steps being performed in sequence and repeatedly.

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At the step 216, the event distribution engine 160 awaits confirmation from each intended recipient 172 that the event message 111 was received by that particular intended recipient 172. When the event distribution engine 160 receives confirmation from all intended recipients 172, the method proceeds with the next step.

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At the step 217, the event distribution engine 160 removes the information about the event message 111 from the event indication queue 150 and from the persistent memory 140.

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At the flow point 220, the system 100 has completely processed the event 6 7 message 111.

#### Replaying Event Messages 2.

At a flow point 230, the system 100 has recovered from a system crash or a system error.

At a step 231, the event replay engine 190 reads information about event messages 111 from the persistent memory 140. As part of this step, the event replay engine 190 performs three sub-steps 231(a), 231(b), and 231(c).

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At the sub-step 231(a), the system initialization replay sub-element 191 reads information about event messages 111 associated with the pre-allocated postinitialization event message resources 130. The event replay engine 190 replays these event messages 111.

At the sub-step 231(b), the incomplete event distribution replay subelement 192 reads information about event messages 111 associated with the preallocated initialization event message resources 120. The event replay engine 190 replays these event messages 111 next.

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At the sub-step 231(c), the cooperating systems replay sub-element 193 reads information about event messages 111, from the persistent memory 140, associated with and stored there by a cooperating system 100. The event replay engine 190 replays these event messages 111 only if the cooperating system 100 is not operational at the time.

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In a preferred embodiment, multiple cooperating systems 100 (preferably a pair of exactly two) are each capable of reading and writing to each other's persistent memories 140. Thus, when a first cooperating system 100 in the pair writes to its persistent memory 140, the second cooperating system 100 in the pair is able to read from that persistent memory 140. If the first cooperating system 100 suffers a system crash or system error, the second cooperating system 100, upon recognizing that system crash or system error, proceeds to replay the event messages 111 from the first cooperating system's persistent memory 140. Operation of multiple cooperating systems 100 is further described in the Incorporated Disclosures, particularly with regard to techniques used to prevent multiple cooperating systems 100 from disrupting each other's operation.

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1	As noted herein, "replay" of event messages 111 is treated by the event in-
2	dication queue 150 and the event distribution engine 160 as if the event messages 111
3	were newly generated. Replayed event messages 111 are processed and delivered before
4	new event messages 111, according to the portion of the method 200 described with re-
5	gard to flow point 210 and flow point 220.
6	
7	At a flow point 240, the system 100 has replayed all event messages 111
8	not yet fully processed, and is ready to proceed at the flow point 210.
<u></u>	
9	3. Multiplexing Recipient Operation
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	A third process in the method 200 is described with regard to a flow point
=12 =13 =14	250, a flow point 260, and steps there-between. Similar to the first process, this third
<u>1</u> 4	process includes a set of three tasks, each of which operates independently and asynchro-
15	nously with regard to each other.
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17	At the flow point 250, a multiplexing recipient 171 is ready to receive ar
18	event message 111.
19	
20	Event Reception

1	A first task includes a sequence including a step 251, a step 252, and a step
2	253. In a preferred embodiment, this first task in its process includes these steps being
3	performed in sequence and repeatedly.
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5	At the step 251, the multiplexing recipient 171 receives the event message
6	111 from the event distribution engine 160.
7	
8	At the step 252, similarly to the steps described with regard to the flow
8 1354 9 4 10 11 11	point 210 and the flow point 220, the multiplexing recipient 171 records information
<b>=</b>	about the event message 111 in its second persistent memory 180.
11	
<b>1</b> 2	At the step 253, the multiplexing recipient 171 (optionally) responds to the
13	event message 111 by confirming that it was received at the multiplexing recipient 171
13 14	(but not necessarily at the intended recipients 172).
15	
16	Event Multiplexing
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18	A second task includes a sequence including a step 254. In a preferred em-
19	bodiment, this second task in its process includes this step being performed repeatedly.
20	
21	At the step 254, the multiplexing recipient 171 (optionally) determines to
22	which intended recipients 172 to deliver the event message 111. In a preferred embodi-

ment, the multiplexing recipient 171 filters the event messages 111 it receives, so that it 1 delivers only those event messages 111 it receives to their actual intended recipients 172. 2 3 For example, a particular intended recipient 172 might determine that it is only interested in a particular subclass of event messages 111. In such cases, the multiplexing recipient 4

171 delivers only that particular subclass of event messages 111 to that particular in-

A third task includes a sequence including a step 255 and a step 256. In a

At the step 255, the multiplexing recipient 171 awaits confirmation from

preferred embodiment, this third task in its process includes these steps being performed

each particular intended recipient 172 that the particular intended recipient 172 has re-

6 tended recipient 172.

in sequence and repeatedly.

ceived the event message 111.

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# **Event Confirmation**

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At the step 256, the multiplexing recipient 171 receives such confirmation from individual intended recipients 172. As part of this step, the multiplexing recipient 171 (optionally) forwards those confirmations on to the event distribution engine 160. When the multiplexing recipient 171 receives all such confirmations, it removes the information about the event message 111 from the second persistent memory 180.

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2	At the flow point 260, the multiplexing recipient 171 has completely proc-
3	essed the event message 111.
4	
5	4. Replaying Multiplexed Event Messages
6	
7	At a flow point 270, the multiplexing recipient 171 has recovered from a
8	system crash or a system error.
5 9 F 110 F 11	
10	At a step 271, the multiplexing recipient 171 reads information about even
<u>1</u> 11	messages 111 from the second persistent memory 180.
□12	
12 13 13 14	At a step 272, the multiplexing recipient 171 replays the event messages
14	111 from the second persistent memory 180.
15	
16	"Replay" of event messages 111 by the multiplexing recipient 171 is simi-
17	lar to replay of event messages as described above with regard to the event indication
18	queue 150 and the event distribution engine 160.
19	
20	At a flow point 280, the multiplexing recipient 171 has replayed all even
21	messages 111 not yet fully processed, and is ready to proceed at the flow point 250.

#### Confirming Event Messages 5.

2

1

As described above, there are steps at which the system 100 or the multi-3 plexing recipient 171 awaits confirmation of the event message 111 from the intended re-4 cipient 172. In a preferred embodiment, confirmation of event messages 111 is per-5 formed by each intended recipient 172 as described below with regard to a flow point 6 7 290, a flow point 300, and steps there-between.

At the flow point 290, the intended recipient 172 is ready to receive an event message 111.

At a step 291, the intended recipient 172 receives an event message 111.

At a step 292, the intended recipient 172 parses the event message 111 and processes the event message 111 according to its own (internal) processing rules for that event message 111.

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At a step 293, the intended recipient 172 generates a confirmation message and sends that confirmation message to the sender of the event message 111.

20

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22

At the flow point 300, the intended recipient 172 has received, processed, and confirmed the event message 111. The sender of the event message 111, upon re-

ceipt of the confirmation message, can regard the event message 111 as completely han-1 dled and can safely delete it. 2 3 Generality of the Invention 4 5 The invention has general applicability to various fields of use, not neces-6 sarily related to the services described above. For example, these fields of use can include 7 one or more of, or some combination of, the following: 9 10 11 11 11 11 11 11 11 The invention is applicable to persistent and reliable delivery of messages other than event messages. The invention is applicable to persistent and reliable operation of other processes than delivery of messages. 15 The invention is applicable to mutually cooperating systems to perform other per-16 sistent and reliable operations. 17 18 The invention is applicable to hierarchical cooperating systems to perform other 19 persistent and reliable operations. 20 21

Other and further applications of the invention in its most general form, will be clear to those skilled in the art after perusal of this application, and are within the scope and spirit of the invention.

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Although preferred embodiments are disclosed herein, many variations are possible which remain within the concept, scope, and spirit of the invention, and these variations would become clear to those skilled in the art after perusal of this application.

1	CLAIMS
2	
3	1. A method including
4	persistently maintaining at least one event message until at least one in
5	tended recipient of said event message confirms delivery of said event message; and
6	upon recovery from an error, replaying said event message;
7	whereby said event message is reliably delivered to said intended recipient.
8	
8 19 10 11	2. A method as in claim 1, including
- - - - 10	receiving said event message by said intended recipient; and
11	generating a confirmation of said event message in response to said even
12	message.
12 13 14	
14	3. A method as in claim 1, wherein said event message is provided by
15	at least one event message producer.
16	
17	4. A method as in claim 1, wherein said persistent maintenance in
18	cludes recording said event message in an event-indication queue, said event-indication
19	queue having resources pre-allocated before occurrence of an event associated with said
20	event message.
21	

1	5. A method as in claim 1, wherein said persistent maintenance in-
2	cludes recording said event message in an event-indication queue, wherein said event-
3	indication queue is reliable even when the event message indicates that allocation of new
4	resources is unstable.
5	
6	6. A method as in claim 1, wherein said persistent maintenance in-
7	cludes recording said event message during a duration when delivery of said event mes-
8	sage is not yet feasible.
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
二月 9 月 10 月 11 月 11	7. A method as in claim 6, including
11	upon termination of said duration, replaying said event message;
<u>-</u> 12	whereby said event message is reliably delivered to said intended recipient.
12 13 13 14	
<b>14</b>	8. A method as in claim 6, wherein said duration includes a boot time
15	or an initialization time.
16	
17	9. A method as in claim 1, wherein said persistent maintenance in-
18	cludes recording said event message in a persistent memory.
19	
20	10. A method as in claim 9, including
21	delivering said event message to said intended recipient;
22	receiving a confirmation of said delivery; and

1	removing said event message from said persistent memory in response to
2	said confirmation.
3	
4	11. A method including
5	persistently maintaining at least one event message during a duration when
6	delivery of said event message is not yet feasible; and
7	upon termination of said duration, replaying said event message;
8	whereby said event message is reliably delivered to an intended recipient of
五 万 9	said event message.
<u> </u>	
8 9 0 1 1 2 1 3 4	12. A method as in claim 11, wherein said duration includes a boot time
12	or an initialization time.
13	
<u>-</u> -14	13. A method as in claim 11, wherein said event message is provided by
15	at least one event message producer.
16	
17	14. A method as in claim 11, including persistently maintaining at least
18	one event message until at least one intended recipient of said event message confirms
19	delivery of said event message.
20	
21	15. A method as in claim 14, including
22	upon recovery from an error, replaying said event message;

1	whereby said event message is reliably delivered to said intended recipient.
2	
3	16. A method as in claim 14, wherein said persistent maintenance in-
4	cludes recording said event message in an event-indication queue, said event-indication
5	queue having resources pre-allocated before occurrence of an event associated with said
6	event message.
7	
8	17. A method as in claim 14, wherein said persistent maintenance in-
5 9 F 110 F 11	cludes recording said event message in an event-indication queue, wherein said event-
10 10	indication queue is reliable even when the event message indicates that allocation of new
11	resources is unstable.
12 13 13 14	
=13	18. A method as in claim 11, wherein said persistent maintenance in-
<u>=</u> 14	cludes recording said event message in a persistent memory.
15	
16	19. A method as in claim 18, including
17	delivering said event message to said intended recipient;
18	receiving a confirmation of said delivery; and
19	removing said event message from said persistent memory in response to
20	said confirmation.
21	
22	20. A method as in claim 11, including

1	receiving said event message by said intended recipient; and
2	generating a confirmation of said event message in response to said event
3	message.
4	
5	21. A method including
6	maintaining at least one event message in a plurality of memory locations,
7	each one of said plurality of memory locations being accessible by both a first server de-
8	vice and a second server device; and
8 19 10 10 11	upon recovery from an error at said first server device, replaying said event
<u>-</u> 110	message from said second server device;
1	whereby said event message is reliably delivered to an intended recipient of
12	said event message.
12 13 14	
<b>=</b> 14	22. A method as in claim 21, wherein said event message is provided by
15	at least one event message producer.
16	
17	23. A method as in claim 21, wherein said maintenance includes persis-
18	tently maintaining said event message during a duration when delivery of said event mes-
19	sage is not yet feasible.
20	
21	24. A method as in claim 23, including
22	upon termination of said duration, replaying said event message;

1	whereby said event message is reliably delivered to an intended recipient of
2	said event message.
3	
4	25. A method as in claim 23, wherein said duration includes a boot time
5	or an initialization time.
6	
7	26. A method as in claim 23, wherein said event message is provided by
8	at least one event message producer.
9	
10	27. A method as in claim 21, wherein said maintenance includes persis-
11	tently maintaining said event message until at least one intended recipient of said event
12	message confirms delivery thereof.
13	
14	28. A method as in claim 27, wherein said persistent maintenance in-
15	cludes recording said event message in an event-indication queue, said event-indication
16	queue having resources pre-allocated before occurrence of an event associated with said
17	event message.
18	
19	29. A method as in claim 27, wherein said persistent maintenance in-
20	cludes recording said event message in an event-indication queue, wherein said event-
21	indication queue is reliable even when the event message indicates that allocation of new
22	resources is unstable.

1	
2	30. A method as in claim 27, wherein said persistent maintenance in-
3	cludes recording said event message in a persistent memory.
4	
5	31. A method as in claim 30, including
6	delivering said event message to said intended recipient;
7	receiving a confirmation of said delivery; and
8	removing said event message from said persistent memory in response to
直 〒 9	said confirmation.
<u>=</u>	
8 9 9 0 1 1 1 1 2 1 3 1 1 4 1 4 1 4 1 4 1 4 1 4 1 4 1 4 1	32. A method as in claim 30, including
<b>_1</b> 2	receiving said event message by said intended recipient; and
⊉ =13 =	generating a confirmation of said event message in response to said event
14	message.
15	
16	33. A method including
17	delivering at least one event message to a multiplexing recipient;
18	maintaining said event message at said multiplexing recipient; and
19	reliably delivering said event message from said multiplexing recipient to at
20	least one intended recipient of said event message.
21	
22	34. A method as in claim 33, including

1	receiving said event message by said intended recipient; and
2	generating a confirmation of said event message in response to said event
3	message.
4	
5	35. A method as in claim 33, wherein said event message is provided by
6	at least one event message producer.
7	
8	36. A method as in claim 33, wherein reliable delivery of said event
3 9 9 10 110 111 111 111 111 111 111 111	message from said multiplexing recipient includes
<u>-</u> 10	persistently maintaining said event message at said multiplexing recipient;
11	upon recovery from an error at said multiplexing recipient, replaying said
<b>12</b>	event message from said multiplexing recipient to said intended recipient;
13	whereby said event message is reliably delivered to said intended recipient.
13 13 14	
15	37. A method as in claim 36, wherein said persistent maintenance in-
16	cludes recording said event message in an event-indication queue, said event-indication
17	queue having resources pre-allocated before occurrence of an event associated with said
18	event message.
19	
20	38. A method as in claim 36, wherein said persistent maintenance in-
21	cludes recording said event message in an event-indication queue, wherein said event-

1	indication queue is reliable even when the event message indicates that allocation of new		
2	resources is unstable.		
3			
4	39. A method as in claim 36, wherein said persistent maintenance in-		
5	cludes recording said event message in a persistent memory.		
6			
7	40. A method as in claim 39, including		
8	delivering said event message to said intended recipient;		
<u>=</u> = 9	receiving a confirmation of said delivery; and		
10	removing said event message from said persistent memory in response to		
8 9 0 1 1 2 2 3 4	said confirmation.		
<u>1</u> 2			
五 十3	41. A method as in claim 33, wherein reliable delivery of said event		
<b>=</b> 14	message from said multiplexing recipient includes		
15	persistently maintaining said event message at said multiplexing recipient		
16	until at least one intended recipient of said event message confirms delivery of said event		
17	message;		
18	sending a confirmation of delivery from said multiplexing recipient in re-		
19	sponse to a confirmation of delivery from said intended recipient.		
20			
21	42. A method as in claim 41, wherein said persistent maintenance in		
22	cludes recording said event message in an event-indication queue, said event-indication		

1	queue having resources pre-allocated before occurrence of an event associated with said
2	event message.
3	
4	43. A method as in claim 41, wherein said persistent maintenance in-
5	cludes recording said event message in an event-indication queue, wherein said event-
6	indication queue is reliable even when the event message indicates that allocation of new
7	resources is unstable.
8 9 10 11 12 13 14	44. A method as in claim 36, wherein said persistent maintenance in-
10 11	cludes recording said event message in a persistent memory.
<u>=</u> 12	45. A method as in claim 44, including
三 三 三 三	delivering said event message to said intended recipient;
<u>=</u> 14	receiving a confirmation of said delivery; and
15	removing said event message from said persistent memory in response to
16	said confirmation.
17	
18	46. A memory including instructions, said instructions capable of being
19	interpreted to indicate
20	persistently maintaining at least one event message until at least one in-
21	tended recipient of said event message confirms delivery of said event message; and
22	upon recovery from an error, replaying said event message;

1	whereby said event message is reliably delivered to said intended recipient.
2	
3	47. A memory as in claim 46, wherein said instructions are also capable
4	of being interpreted to indicate recording said event message during a duration when de-
5	livery of said event message is not yet feasible.
6	
7	48. A memory including instructions, said instructions capable of being
8	interpreted to indicate
9 10 10 11	maintaining at least one event message in a plurality of memory locations,
<b>4</b> 0	each one of said plurality of memory locations being accessible by both a first server de-
1	vice and a second server device; and
<b>1</b> 2	upon recovery from an error at said first server device, replaying said event
13 113	message from said second server device;
13 14	whereby said event message is reliably delivered to an intended recipient of
15	said event message.
16	
17	49. A memory including instructions, said instructions capable of being
18	interpreted to indicate
19	delivering at least one event message to a multiplexing recipient;
20	maintaining said event message at said multiplexing recipient; and
21	reliably delivering said event message from said multiplexing recipient to at
22	least one intended recipient of said event message.

1
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50.	Apparatus	1110	ludina	٦
20.	Abbaratus	HIO.	լապոււ	<

means for persistently maintaining at least one event message until at least one intended recipient of said event message confirms delivery of said event message; and

means for replaying said event message upon recovery from an error.

## 51. Apparatus as in claim 50, including

means for receiving said event message by said intended recipient; and means for generating a confirmation of said event message in response to said event message.

52. Apparatus as in claim 50, wherein said means for persistently maintaining includes means for recording said event message in an event-indication queue, said event-indication queue having resources pre-allocated before occurrence of an event associated with said event message.

53. Apparatus as in claim 50, wherein said means for persistently maintaining includes means for recording said event message in an event-indication queue, wherein said event-indication queue is reliable even when the event message indicates that allocation of new resources is unstable.

1	54. Apparatus as in claim 50, wherein said means for persistently main-		
2	taining includes means for recording said event message during a duration when delivery		
3	of said event message is not yet feasible.		
4			
5	55. Apparatus as in claim 54, including		
6	upon termination of said duration, means for replaying said event message;		
7	whereby said event message is reliably delivered to said intended recipient.		
_ 8			
8 9 4 7 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	56. Apparatus as in claim 54, wherein said duration includes a boot time		
<u>=</u>	or an initialization time.		
<u> </u>			
11	57. Apparatus as in claim 50, wherein said means for persistently main-		
12 13 14	taining includes means for recording said event message in a persistent memory.		
<u>1</u> 4			
15	58. Apparatus as in claim 57, including		
16	means for delivering said event message to said intended recipient;		
17	means for receiving a confirmation of said delivery; and		
18	means for removing said event message from said persistent memory in re-		
19	sponse to said confirmation.		
20	4		
21	59. Apparatus including		

1	means for persistently maintaining at least one event message during a du-
2	ration when delivery of said event message is not yet feasible; and
3	upon termination of said duration, means for replaying said event message.
4	
5	60. Apparatus as in claim 59, wherein said duration includes a boot time
6	or an initialization time.
7	
8	61. Apparatus as in claim 59, including means for persistently main-
<u>-</u>	taining at least one event message until at least one intended recipient of said event mes-
8 9 0 11 2 3 4	sage confirms delivery of said event message.
1	
<b>1</b> 2	62. Apparatus as in claim 61, including, upon recovery from an error
<b>1</b> 3	means for replaying said event message.
<u>=</u> 4	
15	63. Apparatus as in claim 61, wherein said means for persistently main-
16	taining includes means for recording said event message in an event-indication queue
17	said event-indication queue having resources pre-allocated before occurrence of an even
18	associated with said event message.
19	
20	64. Apparatus as in claim 61, wherein said means for persistently main
21	taining includes means for recording said event message in an event-indication queue

1	wherein said event-indication queue is reliable even when the event message indicates
2	that allocation of new resources is unstable.
3	
4	65. Apparatus as in claim 59, wherein said means for persistently main-
5	taining includes means for recording said event message in a persistent memory.
6	
7	66. Apparatus as in claim 65, including
1	means for delivering said event message to said intended recipient;
8 9 4 0	means for receiving a confirmation of said delivery; and
10	means for removing said event message from said persistent memory in re-
<u></u> 1	sponse to said confirmation.
12	
12 113 114	67. Apparatus as in claim 65, including
1 14	means for receiving said event message by said intended recipient; and
15	means for generating a confirmation of said event message in response to
16	said event message.
17	
18	68. Apparatus including
19	means for maintaining at least one event message in a plurality of memory
20	locations, each one of said plurality of memory locations being accessible by both a first
21	server device and a second server device; and

1		upon 1	recovery from an error at said first server device, means for replaying
2	said event m	essage :	from said second server device.
3			
4		69.	Apparatus including
5		means	for delivering at least one event message to a multiplexing recipient;
6		means	for maintaining said event message at said multiplexing recipient;
7	and		
<b>=</b> 8		means	for reliably delivering said event message from said multiplexing re-
19 119	cipient to at l	east on	e intended recipient of said event message.
10			
8 9 1 1 1 1 1 1 1 1 1 1 1 1 1 1		70.	Apparatus as in claim 69, including
<b>-1</b> 2		means	for receiving said event message by said intended recipient; and
13		means	for generating a confirmation of said event message in response to
<u> </u>	said event me	essage.	
15			
16		71.	In a method including reliable delivery of event messages, a memory
17	including		
18		a persi	istent record of at least one event message until at least one intended
19	recipient of s	aid eve	nt message confirms delivery of said event message; and
20		upon r	recovery from an error, a replayable instance of said event message.
21			

1	72. A memory as in claim 71, including a record of said event message
2	during a duration when delivery of said event message is not yet feasible.
3	
4	73. A memory as in claim 71, including
5	at least one event message in a plurality of memory locations, each one of
6	said plurality of memory locations being accessible by both a first server device and a
7	second server device; and
_8	upon recovery from an error at said first server device, at least one instance
₩ 1779 E	of said event message replayable from said second server device.
-8 	
	74. In a method including reliable delivery of event messages, a memory
<b>1</b> 2	including
12 13 14	a persistent record of at least one event message at a multiplexing recipient;
<u>1</u> 4	and
15	an instance of said event message deliverable from said multiplexing re-
16	cipient to at least one intended recipient of said event message.
17	
18	75. In apparatus having elements capable of performing a method, said
19	method including reliable delivery of event messages, a memory including
20	a persistent record of at least one event message until at least one intended
21	recipient of said event message confirms delivery of said event message; and
22	upon recovery from an error, a replayable instance of said event message.

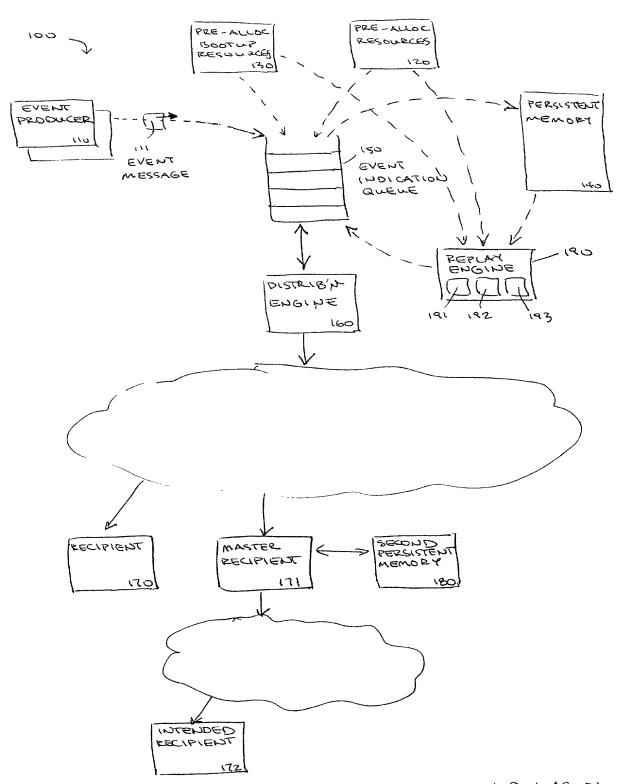
1	
2	76. A memory as in claim 75, including a record of said event message
3	during a duration when delivery of said event message is not yet feasible.
4	
5	77. A memory as in claim 75, including
6	at least one event message in a plurality of memory locations, each one of
7	said plurality of memory locations being accessible by both a first server device and a
<b>8</b>	second server device; and
8 9 0	upon recovery from an error at said first server device, at least one instance
10	of said event message replayable from said second server device.
12 13 13	78. A memory as in claim 75, including
<u>=</u> 13	a persistent record of at least one event message at a multiplexing recipient;
<u> </u>	and
15	an instance of said event message deliverable from said multiplexing re-

cipient to at least one intended recipient of said event message.

16

### ABSTRACT OF THE DISCLOSURE

The invention provides a method and system for persistent and reliable de-
livery of event messages. Those parts of the system responsible for delivering event
messages are able to persistently maintain those event messages until the intended recipi-
ent of the event message confirms delivery of those event messages. Those parts of the
system responsible for recovering from system crashes and other system errors are able to
persistently maintain those event messages until delivery, even after recovery from sys-
tem crashes or other system errors. The system includes a set of event message produc-
ers, and maintains an event-indication queue of those event messages provided by the
event producers using a set of pre-allocated resources. An event-distribution engine dis-
tributes event messages to intended recipients and, after having received confirmation
that the event messages were received, removes them from the event-indication queue.
Recipients of event messages receive the event messages, acknowledge their receipt
thereof, and might take action in response to the event message. The system includes
persistent memory, initialization memory, and recipient persistent memories, and pro-
vides upon recovery from system crashes or other system error, an ability to replay event
messages recorded in those memories, to re-present them as event messages. A cluster of
file servers collectively forming a highly-available system shares persistent memories, so
that upon a system crash or other system error, at least one other file server has a record
of those event messages.



103-1048-01 FIGURE 1. SHEET 1 Of 4

#### 1. Processing Event Messages

210

The system 100 is ready to receive an event message 111.

211

An event producer 110 generates an event message 111.

212

If the system 100 is in normal operation, the event message 111 is recorded in the pre-allocated initialization event message resources 120. If the system 100 is still in its initialization time duration, the event message 111 is recorded in the preallocated post-initialization message resources 130.

213

The system 100 copies information about the event message 111 to a set of locations in the persistent memory 140.

214

The system 100 copies information about the event message 111 to the event indication queue 150. In a preferred embodiment, the event indication queue 150 includes a FIFO similar to that maintained in the persistent memory 140.

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2. Replaying Event Messages

230 The system 100 has recovered from a system crash or a system error.

to Fig 2B

215

The event distribution engine 160 responds to the information about the event message 111 in the event indication queue 150. The event distribution engine 160 delivers the event message 111 to its intended recipients 172.

216

The event distribution engine 160 confirmation awaits from each intended recipient 172 that the event message 111 was received by that particular intended recipient 172. When the event distribution engine 160 receives confirmation from all intended recipients 172, the method proceeds with the next step.

217

The event distribution engine 160 removes the information about the event message 111 from the event indication queue 150 and from the persistent memory 140.

<u> 220</u>

The system 100 has completely processed the event message 111.

from Fig 2A

The event replay engine 190 reads information about event messages 111 from the persistent memory 140.

**√**231(a)

The system initialization replay subelement 191 reads information about event messages 111 associated with the pre-allocated post-initialization event message resources 130. The event replay engine 190 replays these event messages 111.

231(b)

The incomplete event distribution replay sub-element 192 reads information about event messages 111 associated with the pre-allocated initialization event message resources 120. The event replay engine 190 replays these event messages 111 next.

231(c)

The cooperating systems replay subelement 193 reads information about event messages 111, from the persistent memory 140, associated with and stored there by a cooperating system 100.

240

The system 100 has replayed all event messages 111 not yet fully processed, and is ready to proceed at the flow point 210.

3. Multiplexing Recipient Operation

250

A multiplexing recipient 171 is ready to receive an event message 111.

1

251

The multiplexing recipient 171 receives the event message 111 from the event distribution engine 160.

252

The multiplexing recipient 171 records information about the event message 111 in its second persistent memory 180.

253

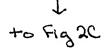
The multiplexing recipient 171 (optionally) responds to the event message 111 by confirming that it was received at the multiplexing recipient 171 (but not necessarily at the intended recipients 172).

254

The multiplexing recipient 171 (optionally) determines to which intended recipients 172 to deliver the event message 111.

<u> 255</u>

The multiplexing recipient 171 awaits confirmation from each particular intended recipient 172 that the particular intended recipient 172 has received the event message 111.



From 2B

The multiplexing recipient 171 receives such confirmation from individual intended recipients 172.

260

The multiplexing recipient 171 has completely processed the event message 111.

# 4. Replaying Multiplexed Event Messages

270

The multiplexing recipient 171 has recovered from a system crash or a system error.

271

The multiplexing recipient 171 reads information about event messages 111 from the second persistent memory 180.

272

The multiplexing recipient 171 replays the event messages 111 from the second persistent memory 180.

280

The multiplexing recipient 171 has replayed all event messages 111 not yet fully processed, and is ready to proceed at the flow point 250.

#### 5. Confirming Event Messages

290

The intended recipient 172 is ready to receive an event message 111.

291

The intended recipient 172 receives an event message 111.

292

The intended recipient 172 parses the event message 111 and processes the event message 111 according to its own (internal) processing rules for that event message 111.

<u> 293</u>

The intended recipient 172 generates a confirmation message and sends that confirmation message to the sender of the event message 111.

300

The intended recipient 172 has received, processed, and confirmed the event message 111. The sender of the event message 111, upon receipt of the confirmation message, can regard the event message 111 as completely handled and can safely delete it.